

THISTLETHWAITE'S THEOREM FOR VIRTUAL LINKS

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ABSTRACT. The celebrated Thistlethwaite theorem relates the Jones polynomial of a link with the Tutte polynomial of the corresponding planar graph. We give a generalization of this theorem to virtual links. In this case, the graph will be embedded into a (higher genus) surface. For such graphs we use the generalization of the Tutte polynomial discovered by B. Bollobás and O. Riordan.

INTRODUCTION

Regions of a link diagram can be colored black and white in a checkerboard pattern. Putting a vertex in each black region and connecting two vertices by an edge if the corresponding regions share a crossing yields a planar graph. In 1987 Thistlethwaite [Th] proved that the Jones polynomial of an alternating link can be obtained as a specialization of the Tutte polynomial of the corresponding planar graph. L. Kauffman [K2] generalized the theorem to arbitrary links using signed graphs and extending the Tutte polynomial to them. An expression for the Jones polynomial in terms of the Bollobás-Riordan polynomial, without signed graphs, was found in [DFKLS]. The idea to use the Bollobás-Riordan polynomial instead of the Tutte polynomial belongs to Igor Pak. It was first realized in [CP], where Thistlethwaite's theorem was generalized to checkerboard colorable virtual links.

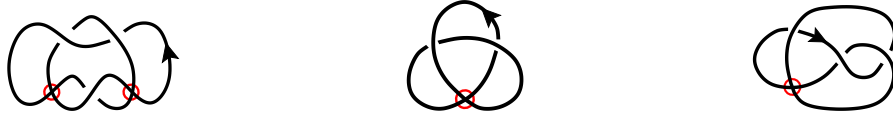
Here we shall generalize this theorem to *arbitrary* virtual links.

We recall the basic definitions of virtual links, their Jones polynomial through the Kauffman bracket, ribbon graphs, and the Bollobás-Riordan polynomial in Sections 1 and 2. The key construction of a ribbon graph from a digram of a virtual link is explained in Section 3. Our main theorem is formulated and proved in Section 4.

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1. VIRTUAL LINKS AND THE KAUFFMAN BRACKET

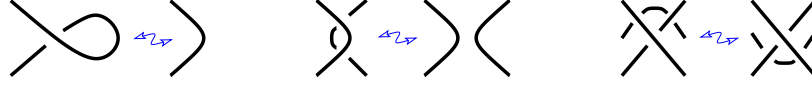
The theory of *virtual links* was discovered independently by L. Kauffman [K3] and M. Gousarov, M. Polyak, and O. Viro [GPV] around 1998. According to Kauffman's approach, virtual links are represented by diagrams similar to ordinary knot diagrams, except some crossings are designated as *virtual*. Virtual crossings should be understood not as crossings but rather as defects of our two-dimensional pictures. They should be treated in the same way as the extra crossings appearing in planar pictures of non-planar graphs. Here are some examples.



The virtual crossings in these pictures are circled to distinguish them from the classical ones.

Key words and phrases. Virtual knots and links, ribbon graph, Kauffman bracket, Bollobás-Riordan polynomial.

Virtual link diagrams are considered up to the *classical* Reidemeister moves involving classical crossings:



and the *virtual* Reidemeister moves:



The Kauffman bracket and the Jones polynomial for virtual links are defined in the same way as for classical ones. Let L be a virtual link diagram. Consider two ways of resolving a classical crossing. The *A-splitting*, $\frac{\text{crossing}}{\text{arc}} \rightsquigarrow \text{arc}$, is obtained by joining the two vertical angles swept out by the overcrossing arc when it is rotated counterclockwise toward the undercrossing arc. Similarly, the *B-splitting*, $\frac{\text{crossing}}{\text{arc}} \rightsquigarrow \text{arc}$, is obtained by joining the other two vertical angles. A *state* S of a link diagram L is a choice of either an *A*- or *B*-splitting at each classical crossing of the diagram. Denote by $\mathcal{S}(L)$ the set of the states of L . Clearly, a diagram L with n crossings has $|\mathcal{S}(L)| = 2^n$ different states.

Denote by $\alpha(S)$ and $\beta(S)$ the numbers of *A*-splittings and *B*-splittings in a state S , respectively. Also, denote by $\delta(S)$ the number of components of the curve obtained from the link diagram L by splitting according to the state $S \in \mathcal{S}(L)$. Note that virtual crossings do not connect components.

Definition 1.1. The Kauffman bracket of a diagram L is a polynomial in three variables A , B , d defined by the formula:

$$[L](A, B, d) := \sum_{S \in \mathcal{S}(L)} A^{\alpha(S)} B^{\beta(S)} d^{\delta(S)-1}.$$

Note that $[L]$ is *not* a topological invariant of the link and in fact depends on the link diagram. However, it defines the *Jones polynomial* $J_L(t)$ by a simple substitution:

$$J_L(t) := (-1)^{w(L)} t^{3w(L)/4} [L](t^{-1/4}, t^{1/4}, -t^{1/2} - t^{-1/2}).$$

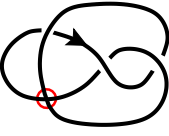
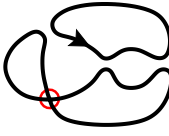
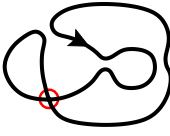
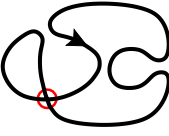
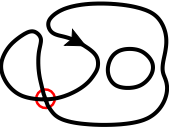
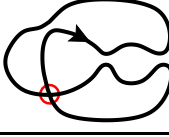
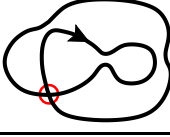
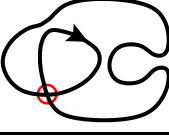
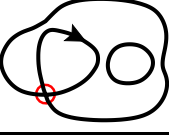
Here $w(L)$ denotes the *writhe*, determined by orienting L and taking the sum over the classical crossings of L of the following signs:



The Jones polynomial is a classical topological invariant (see e.g. [K1]).

Example 1.2. Consider the third virtual knot diagram L from the example above. It is shown on the left of the table below. It has one virtual and three classical crossings (one positive and two negative). So there are eight states and $w(L) = -1$. The curves obtained by the splittings and the corresponding parameters $\alpha(S)$, $\beta(S)$, and $\delta(S)$ are shown in the remaining columns of

the table.

				
(α, β, δ)	$(3, 0, 2)$	$(2, 1, 1)$	$(2, 1, 1)$	$(1, 2, 2)$
				
	$(2, 1, 1)$	$(1, 2, 1)$	$(1, 2, 1)$	$(0, 3, 2)$

We have

$$\begin{aligned}
 [L] &= A^3 d + 3A^2 B + 2AB^2 + AB^2 d + B^3 d ; \\
 J_L(t) &= t^{-2} - t^{-1} - t^{-1/2} + 1 + t^{1/2} .
 \end{aligned}$$

2. RIBBON GRAPHS AND THE BOLLOBÁS-RIORDAN POLYNOMIAL

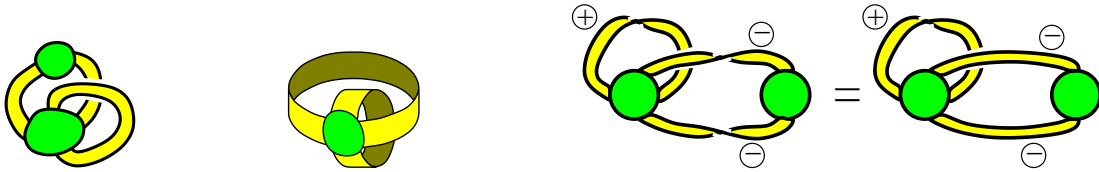
Ribbon graphs are the objects of Topological Graph Theory. There are several books on this subject and its applications [GT, LZ, MT]. Our ribbon graphs are nothing else than the band decompositions from [GT, Section 3.2] with the interior of all 2-bands removed. We modify the definitions of [BR] to signed ribbon graphs.

Definition 2.1. A (signed) *ribbon graph* G is a surface (possibly nonorientable) with boundary represented as the union of two sets of closed topological discs called *vertices* $V(G)$ and *edges* $E(G)$, satisfying the following conditions:

- these vertices and edges intersect by disjoint line segments;
- each such line segment lies on the boundary of precisely one vertex and precisely one edge;
- every edge contains exactly two such line segments,

together with a *sign function* $\varepsilon : E(G) \rightarrow \{\pm 1\}$.

Here are a few examples (if the sign is omitted it is assumed to be +1).



If we put a dot at the center of each vertex-disc and a line in each edge-disc we will get an ordinary graph Γ , the *core graph*, embedded into a surface of G . Conversely if we have a graph Γ embedded into a surface then it determines a ribbon graph structure on a small neighborhood of Γ inside the surface.

To define the Bollobás-Riordan polynomial we need to introduce several parameters of a ribbon graph G . Let

- $v(G) := |V(G)|$ denote the number of vertices of G ;
- $e(G) := |E(G)|$ denote the number of edges of G ;
- $k(G)$ denote the number of connected components of G ;
- $r(G) := v(G) - k(G)$ be the *rank* of G ;
- $n(G) := e(G) - r(G)$ be the *nullity* of G ;

- $\text{bc}(G)$ denote the number of connected components of the boundary of the surface of G .

A *spanning subgraph* of a ribbon graph G is defined as a subgraph consisting of all the vertices of G and a subset of the edges of G . Let $\mathcal{F}(G)$ denote the set of spanning subgraphs of G . Clearly, $|\mathcal{F}(G)| = 2^{e(G)}$. For a signed ribbon graph we need one more parameter of a spanning subgraph. Let $e_-(F)$ be the number of negative edges in F . Denote the complement to F in G by $\bar{F} = G - F$, i.e. the spanning subgraph of G with exactly those (signed) edges of G that do not belong to F . Finally, let

$$s(F) = \frac{e_-(F) - e_-(\bar{F})}{2}.$$

Definition 2.2. The signed *Bollobás-Riordan polynomial* $R_G(x, y, z)$ is defined by

$$R_G(x, y, z) := \sum_{F \in \mathcal{F}(G)} x^{r(G) - r(F) + s(F)} y^{n(F) - s(F)} z^{k(F) - \text{bc}(F) + n(F)}.$$

In general this is a Laurent polynomial in $x^{1/2}$, $y^{1/2}$, and z .

The signed version of the Bollobás-Riordan polynomial was introduced in [CP]. If all the edges are positive then it is obtained from the original Bollobás-Riordan polynomial [BR] by a simple substitution $x + 1$ for x . Note that the exponent $k(F) - \text{bc}(F) + n(F)$ of the variable z is equal to $2k(F) - \chi(\tilde{F})$, where $\chi(\tilde{F})$ is the Euler characteristic of the surface \tilde{F} obtained by gluing a disc to each boundary component of F . For orientable ribbon graphs it is twice the genus of F . In particular, for a planar ribbon graph G (i.e. when the surface G has genus zero) the Bollobás-Riordan polynomial R_G does not depend on z . In this case it is essentially equal to the classical Tutte polynomial $T_\Gamma(x, y)$ of the core graph Γ of G :

$$R_G(x - 1, y - 1, z) = T_\Gamma(x, y)$$

if all edges are positive, and if not, to Kauffman's signed Tutte polynomial for signed graphs. Similarly, a specialization $z = 1$ of the Bollobás-Riordan polynomial of an arbitrary ribbon graph G gives the (signed) Tutte polynomial of the core graph:

$$R_G(x - 1, y - 1, 1) = T_\Gamma(x, y).$$

So one may think about the Bollobás-Riordan polynomial as a generalization of the Tutte polynomial to graphs embedded into a surface.

Example 2.3. Consider the third ribbon graph G from our example above and shown on the left in the table below. The other columns show eight possible spanning subgraphs F and the corresponding values of $k(F)$, $r(F)$, $n(F)$, $\text{bc}(F)$, and $s(F)$.

(k, r, n, bc, s)	$(1, 1, 1, 2, 1)$	$(1, 1, 0, 1, 0)$	$(1, 1, 0, 1, 0)$	$(2, 0, 0, 2, -1)$
	$(1, 1, 2, 1, 1)$	$(1, 1, 1, 1, 0)$	$(1, 1, 1, 1, 0)$	$(2, 0, 1, 2, -1)$

We have

$$R_G(x, y, z) = x + 2 + y + xyz^2 + 2yz + y^2z.$$

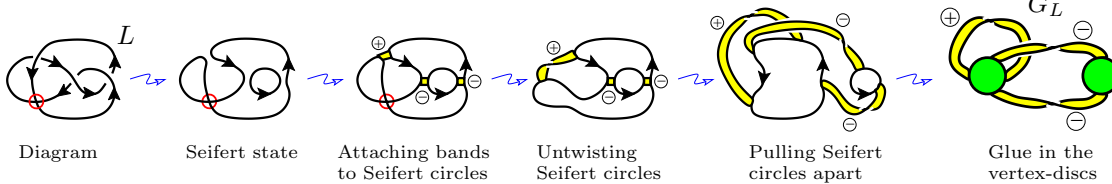
3. RIBBON GRAPHS ASSOCIATED WITH VIRTUAL LINKS

In this section we describe a construction of a ribbon graph starting with a virtual link diagram. Our construction is similar to the classical Seifert algorithm of a construction of the Seifert surface of a link. There are two differences. The first one is that we do not twist the bands in small neighborhoods of the crossings. The second one is that we do not care how our ribbon graph is embedded into the three space.

Suppose our virtual link diagram L is oriented. Then there is a state where all the splittings preserve orientation. Following a suggestion of N. Stoltzfus, we will call it the *Seifert state*, because its state circles are the Seifert circles of the link diagram. Also we will call all splittings in the Seifert state *Seifert splittings*.

The Seifert circles will be the boundary circles of the vertex-discs of the future ribbon graph. So we are going to glue in a disc to each Seifert circle. Before that, though, let us describe the edges of the ribbon graph. When we are doing a Seifert splitting in a vicinity of a crossing we place a small planar band connecting two branches of the splitting. These bands will be the edge-discs of our ribbon graph. If the Seifert splitting was an A -splitting we assign $+1$ to the corresponding edge-band, if it was a B -splitting then we assign -1 . It is easy to see that this sign is equal to the local writhe of the crossing. So we get a sign function. Because of the presence of virtual crossings our Seifert circles may be twisted, i.e. they are actually immersed into the plane with double points at virtual crossings. In the next step of the construction we untwist all Seifert circles to resolve the double points. This may result in some twisting on the edge-bands. After that we pull all Seifert circles apart, which could lead to additional twisting of our edges. In the last step, we glue the vertex-discs into the circles. The signed ribbon graph produced is denoted by G_L . The next example illustrates this procedure.

Example 3.1.



Another way to explain the same construction is the following. Instead of attaching bands to the Seifert circles we only mark the places on the Seifert circles where the bands have to be attached and memorize the order in which the marks occur according to the orientation on the circles. Then we draw each Seifert circle separately on a plane as a perfect circle oriented counterclockwise. Now attach edge-bands according to the marks; it is easy to see that this will always result in a half-twist on each band. The sign function is defined as before.

4. MAIN THEOREM

Theorem 4.1. *Let L be a virtual link diagram, G_L be the corresponding signed ribbon graph, and $n := n(G_L)$, $r := r(G_L)$, $k := k(G_L)$. Then*

$$[L](A, B, d) = A^n B^r d^{k-1} R_{G_L} \left(\frac{Ad}{B}, \frac{Bd}{A}, \frac{1}{d} \right).$$

Proof. Let L be a virtual link diagram, G_L be the corresponding signed ribbon graph, and denote the Seifert state of G_L as \bar{S} . There is a natural bijection between $\mathcal{S}(L)$, the set of states of L , and $\mathcal{F}(G_L)$, the set of spanning subgraphs of G_L . Namely, given a state S , associate to it a spanning subgraph F_S by the following construction. If a crossing in S is split differently than it is in \bar{S} , include its associated edge-band in the spanning subgraph F_S . If a crossing is split the same way for both S and \bar{S} , do not include the associated edge-band in F_S . This gives

the subgraph F_S associated to S . (Certainly $|\mathcal{S}(L)| = |\mathcal{F}(G_L)|$, because the number of classical crossings of L is equal to $e(G_L)$ by virtue of our construction above.)

For example, consider the link L from example 1.2. We know from above that G_L is the ribbon graph considered in example 2.3. For each state given in the table on page 3, we can associate to it a spanning subgraph from the table on page 4. Consider the first state S given in the table on page 3. The two rightmost crossings are split differently than they are in the Seifert state of L given above. Thus the spanning subgraph associated to this state S via the correspondence given above is the first subgraph in the table on page 4. In fact, each state in the table on page 3 corresponds correctly to its associated spanning subgraph in the table on page 4. Check that $[L]$ computed in example 1.2 and $R_{G_L}(x, y, z)$ computed in example 2.3 satisfy the theorem.

Now, given that $F \in \mathcal{F}(G_L)$ is associated to $S \in \mathcal{S}(L)$ as described above (for simplicity, we write F instead of F_S), consider the term

$$x^{r(G_L)-r(F)+s(F)} y^{n(F)-s(F)} z^{k(F)-\text{bc}(F)+n(F)} .$$

Substituting in $x = \frac{Ad}{B}$, $y = \frac{Bd}{A}$, and $z = \frac{1}{d}$ and multiplying by the term $A^n B^r d^{k-1}$ as in the theorem, we have

$$\begin{aligned} & A^n B^r d^{k-1} (AdB^{-1})^{r-r(F)+s(F)} (BdA^{-1})^{n(F)-s(F)} d^{-k(F)+\text{bc}(F)-n(F)} \\ &= A^{n+r-r(F)-n(F)+2s(F)} B^{r(F)+n(F)-2s(F)} d^{k+r-k(F)-r(F)+\text{bc}(F)-1} . \end{aligned}$$

Since $r(G) := v(G) - k(G)$ and $n(G) := e(G) - r(G)$ for any ribbon graph G , we can rewrite our term as

$$A^{e(G_L)-e(F)+2s(F)} B^{e(F)-2s(F)} d^{v(G_L)-v(F)+\text{bc}(F)-1} .$$

And since $v(G_L) = v(F)$ by the definition of a spanning subgraph, we have

$$(1) \quad A^{e(G_L)-e(F)+2s(F)} B^{e(F)-2s(F)} d^{\text{bc}(F)-1} .$$

It suffices to show that this is equal to the Kauffman bracket term $A^{\alpha(S)} B^{\beta(S)} d^{\delta(S)-1}$, since our bijection described above will then imply the theorem. We first show that $e(F) - 2s(F) = \beta(S)$ by a counting argument. Using the definition of $s(F)$, we get

$$(2) \quad e(F) - 2s(F) = e(F) - e_-(F) + e_-(\overline{F}) .$$

Consider the crossings of L and how they are split in \overline{S} . Let m denote the number of crossings which are B -splittings in \overline{S} . Let b denote the number of crossings which are B -splittings in \overline{S} but are A -splittings in S . Let a denote the number of crossings which are A -splittings in \overline{S} but are B -splittings in S .

Now, since $e(F)$ is the number of edges included in F , $e(F)$ equals the number of crossings of L which are split differently between S and \overline{S} . That is, $e(F) = a + b$. Recall that $e_-(F)$ denotes the number of edges in F with sign -1 . And since any such edge corresponds to a B -splitting in \overline{S} , it is clear that $e_-(F) = b$.

Since \overline{F} is the complement of F , $e_-(\overline{F})$ denotes the number of crossings which are B -splittings in \overline{S} and also in S . So, we deduce that $e_-(\overline{F}) = m - b$. Finally, we consider $\beta(S)$, the number of crossings which are B -splittings in S . So clearly, $\beta(S) = a + e_-(\overline{F}) = a + (m - b)$.

Thus, we have that

$$\begin{aligned} e(F) - e_-(F) + e_-(\overline{F}) &= (a + b) - b + (m - b) \\ &= a + (m - b) \\ &= \beta(S) . \end{aligned}$$

This with (2) gives the desired result. And the fact that $e(G_L) - e(F) + 2s(F) = \alpha(S)$ is immediate, since we just showed that $e(F) - 2s(F) = \beta(S)$, and certainly $e(G_L) = \alpha(S) + \beta(S)$.

To finish the proof it remains to show that $\text{bc}(F) = \delta(S)$. For that let us trace simultaneously a circle of the state S and a boundary component of F . Suppose we are passing a place near a crossing. If this crossing is split in the same way as in \bar{S} then we continue to go along the circle which is locally the same as in \bar{S} . On the whole ribbon graph G_L this means that we are passing a place where the corresponding edge-band is supposed to be attached. Or, in other words, we skip the edge and do not include it into the corresponding spanning subgraph. This is precisely how we obtained F . Now if the passed crossing in S is split differently compared with \bar{S} , then we have to switch to another strand of \bar{S} . This means that we should turn on the corresponding edge-band in G_L , i.e. we should include this edge-band in the corresponding subgraph. Again this is precisely what we did with F . The next picture with the first state of example 1.2 (the table on page 3) illustrates this.



Therefore, the tracing of the state circles of S corresponds to the tracing of the boundary components of F , i.e. $\text{bc}(F) = \delta(S)$.

So, we have shown that (1) is equal to the term of $[L]$ corresponding to the state S , and thus theorem 4.1 is proved. \square

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